

Some open problems in matchings in graphs

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- Matchings in graphs
- Number of k -matchings in bipartite graphs as permanents and hafnians
- Upper bounds on permanents and hafnians: results and conjectures.
- Lower bounds on permanents and hafnians: results and conjectures.
- Probabilistic methods

Matchings

- $G = (V, E)$ undirected graph with vertices V , edges E .
- matching in G : $M \subseteq E$
no two edges in M share a common endpoint.
- $e = (u, v) \in M$ is dimer
- v not covered by M is monomer.
- M called monomer-dimer cover of G .
- M is perfect matching \iff no monomers.
- M is k -matching $\iff \#M = k$.

Generating matching polynomial

- $\phi(k, G)$ number of k -matchings in G , $\phi(0, G) := 1$
- $\Phi_G(x) := \sum_k \phi(k, G)x^k$ matching generating polyn.
- roots of $\Phi_G(x)$ are real nonpositive Heilmann-Lieb 1972.
Newton inequalities hold
- $\Phi_{G_1 \cup G_2}(x) = \Phi_{G_1}(x)\Phi_{G_2}(x)$

Examples:

$$\Phi_{K_{2r}}(x) = \sum_{k=0}^r \binom{2r}{2k} \frac{\prod_{j=0}^{k-1} \binom{2k-2j}{2}}{k!} x^k = \sum_{k=0}^r \frac{(2r)!}{(2r-2k)!2^k k!} x^k$$

$$\Phi_{K_{r,r}}(x) = \sum_{k=0}^r \binom{r}{k}^2 k! x^k$$

$\mathcal{G}(r, 2n) \supset \mathcal{GB}(r, 2n)$ set of r -regular and regular bipartite graphs on $2n$ vertices, respectively

$qK_{r,r} \in \mathcal{GB}(r, 2rq)$ a union of q copies of $K_{r,r}$.

$$\Phi_{qK_{r,r}} = \Phi_{K_{r,r}}^q$$

Formulas for k -matchings in bipartite graphs

$G = (V, E)$ **bipartite** $V = V_1 \cup V_2, E \subset V_1 \times V_2,$

represented by bipartite adjacency matrix

$$B(G) = B = [b_{ij}]_{i,j=1}^{m \times n} \in \{0, 1\}^{m \times n}, \#V_1 = m, V_2 = n.$$

Example: Any subgraph of \mathbb{Z}^d is bipartite

CLAIM: $\phi(k, G) = \text{perm}_k(B(G)).$

Prf: Suppose $n = \#V_1 = \#V_2.$

Then permutation $\sigma : \langle n \rangle \rightarrow \langle n \rangle$ is a perfect match iff $\prod_{i=1}^n b_{i\sigma(i)} = 1.$

The number of perfect matchings in G is $\phi(n, G) = \text{perm } B(G).$ □

Computing $\phi(n, G)$ is $\#P$ -complete problem Valiant 1979

For $G = (\langle 2n \rangle, E)$ bipartite $G \in \mathcal{GB}(r, 2n) \iff \frac{1}{r}B(G) \in \Omega_n \iff$
 G is a disjoint (edge) union of r perfect matchings

Matching on nonbipartite graphs

$$G = (V, E), |V| = 2n,$$

$$A(G) = [a_{ij}] \in S_0(2n, \{0, 1\}) - \text{adjacency matrix of } G$$

$$\phi(n, G) = \text{haf}(A(G)) = \sum_{M \in \mathcal{M}(K_{2n})} \prod_{(i,j) \in M} a_{ij}$$

$\mathcal{M}(K_{2n})$ the set of perfect matchings in K_{2n}

$$\phi(k, G) = \text{haf}_k(A(G)) = \sum_{M \in \mathcal{M}_k(K_{2n})} \prod_{(i,j) \in M} a_{ij}$$

$\mathcal{M}_k(K_{2n})$ the set of k matchings in K_{2n}

Claim $\text{perm}(A(G)) \geq \text{haf}(A(G))^2$. Equality holds if G is bipartite.

Main problems

Find good estimates on

$$s_n(k, r) := \min_{G \in \mathcal{G}(r, 2n)} \phi(k, G) \leq t_n(k, r) := \min_{G \in \mathcal{GB}(r, 2n)} \phi(k, G)$$

$$S_n(k, r) := \max_{G \in \mathcal{G}(r, 2n)} \phi(k, G) \geq T_n(k, r) := \max_{G \in \mathcal{GB}(r, 2n)} \phi(k, G)$$

Completely solved case $r = 2$ [7]

$S_n(k, 2) = T_n(k, 2)$ achieved only for $G = nK_{2,2}$ or $G = mK_{2,2} \cup C_6$.

$t_n(k, 2)$ achieved only for C_{2n}

$s_n(k, 2)$ achieved only for mC_3 , $mC_3 \cup C_4$ or $mC_3 \cup C_5$.

The upper bound conjecture

$$S_{qr}(k, r) = T_{qr}(k, r) = \phi(k, qK_{r,r})$$

$k = qr$ Follows from Bregman's inequality (see also [3])

$$\text{perm } A \leq \prod_{i=1}^n (r_i!)^{\frac{1}{r_i}}$$

$$A = [a_{ij}] \in \{0, 1\}^{n \times n} \quad r_i = \sum_{j=1}^n a_{ij}, i = 1, \dots, n$$

Egorichev-Alon-Friedland for $G = (V, E), |V| = 2n$

$$\phi(n, G) \leq \prod_{v \in V} (\deg(v)!)^{\frac{1}{2 \deg(v)}}$$

Equality holds iff G a union of complete bipartite graphs

$$S_n(k, r) \leq \binom{2n}{2k} (r!)^{\frac{k}{r}}$$

$$T_n(k, r) \leq \min\left(\binom{n}{k}^2 (r!)^{\frac{k}{r}}, \binom{n}{k} r^k\right)$$

Friedland-Krop-Lundow-Markström [6]

Asymptotic versions

$$Sa(p, r) = \limsup_{n_j \rightarrow \infty, \frac{k_j}{n_j} \rightarrow p \in [0, 1]} \frac{\log S_{n_j}(k_j, r)}{2n_j}$$

$$Ta(p, r) = \limsup_{n_j \rightarrow \infty, \frac{k_j}{n_j} \rightarrow p \in [0, 1]} \frac{\log T_{n_j}(k_j, r)}{2n_j}$$

$$sa(p, r) = \liminf_{n_j \rightarrow \infty, \frac{k_j}{n_j} \rightarrow p \in [0, 1]} \frac{\log s_{n_j}(k_j, r)}{2n_j}$$

$$ta(p, r) = \liminf_{n_j \rightarrow \infty, \frac{k_j}{n_j} \rightarrow p \in [0, 1]} \frac{\log t_{n_j}(k_j, r)}{2n_j}$$

Next slide gives the graphs of AUMC and the upper bounds for $Ta(p, 4)$.

$r = 4$ upper bounds

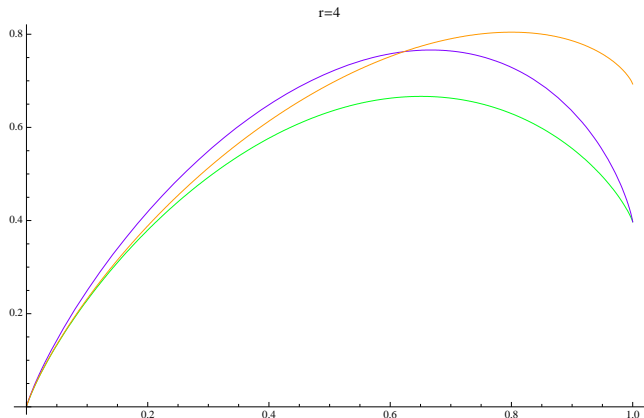


Figure: $h_{K(4)}$ -green, $upp_{4,1}$ -blue, $upp_{4,2}$ -orange

The lower bounds: Bipartite case

$r^k \min_{C \in \Omega_n} \text{perm}_k C \leq \phi(k, G)$ for any $G \in \mathcal{GB}(r, 2n)$

$J_n = B(K_{n,n}) = [1]$ the incidence matrix of the complete bipartite graph $K_{n,n}$ on $2n$ vertices

van der Waerden permanent conjecture 1926:

$$\min_{C \in \Omega_n} \text{perm } C = \text{perm } \frac{1}{n} J_n \left(= \frac{n!}{n^n} \approx \sqrt{2\pi n} e^{-n} \right)$$

Tverberg permanent conjecture 1963:

$$\min_{C \in \Omega_n} \text{perm}_k C = \text{perm}_k \frac{1}{n} J_n \left(= \binom{n}{k}^2 \frac{k!}{n^k} \right)$$

for all $k = 1, \dots, n$.

History

- In 1979 Friedland showed the lower bound $\text{perm } C \geq e^{-n}$ for any $C \in \Omega_n$ following T. Bang's announcement 1976.
This settled the conjecture of Erdős-Rényi on the exponential growth of the number of perfect matchings in $d \geq 3$ -regular bipartite graphs 1968, Voorhoeve 1979.
- van der Waerden permanent conjecture was proved by Egorichev and Falikman 1981.
- Tverberg conjecture was proved by Friedland 1982
- 79 proof is tour de force according to Bang
- 81 proofs involve directly (Egorichev) and indirectly (Falikman) use of Alexandroff mixed volume inequalities with the conditions for the extremal matrix
- 82 proof uses methods of 81 proofs with extra ingredients
- There are new simple proofs using nonnegative hyperbolic polynomials e.g. Friedland-Gurvits 2008

Lower matching bounds for bipartite graphs

Voorhoeve-1979 ($r = 3$) Schrijver-1998

$$\phi(n, \mathbf{G}) \geq \left(\frac{(r-1)^{r-1}}{r^{r-2}}\right)^n \quad \text{for } \mathbf{G} \in \mathcal{GB}(r, 2n)$$

Gurvits 2006: $A \in \Omega_n$, each column has at most r nonzero entries:

$$\text{perm } A \geq \frac{r!}{r^r} \left(\frac{r}{r-1}\right)^{r(r-1)} \left(\frac{r-1}{r}\right)^{(r-1)n}.$$

Cor : $\phi(n, \mathbf{G}) \geq \frac{r!}{r^r} \left(\frac{r}{r-1}\right)^{r(r-1)} \left(\frac{(r-1)^{r-1}}{r^{r-2}}\right)^n$

Con FKM 2006 : $\phi(k, \mathbf{G}) \geq \binom{n}{k}^2 \left(\frac{nr-k}{nr}\right)^{nr-k} \left(\frac{kr}{n}\right)^k, \mathbf{G} \in \mathcal{GB}(r, 2n)$

F-G 2008 showed weaker inequalities

Thm: $r \geq 3, s \geq 1$ integers,

$B_n \in \Omega_n, n = 1, 2, \dots$ each column of B_n has at most r -nonzero entries.

$k_n \in [0, n] \cap \mathbb{N}, n = 1, 2, \dots, \lim_{n \rightarrow \infty} \frac{k_n}{n} = p \in (0, 1]$ then

$$\liminf_{n \rightarrow \infty} \frac{\log \text{perm}_{k_n} B_n}{2n} \geq \frac{1}{2} (-p \log p - 2(1-p) \log(1-p)) + \frac{1}{2} (r+s-1) \log\left(1 - \frac{1}{r+s}\right) - \frac{1}{2} (s-1+p) \log\left(1 - \frac{1-p}{s}\right)$$

Prf combines properties positive hyperbolic polynomials, capacity and the measure on $\mathcal{G}(r, 2n)$

- **Cor:** r -ALMC holds for $p_s = \frac{r}{r+s}, s = 0, 1, \dots,$
- **Con:** under Thm assumptions

$$\liminf_{n \rightarrow \infty} \frac{\log \text{perm}_{k_n} B_n}{2n} \geq f(r, p) - \frac{p}{2} \log r$$

- For $p_s = \frac{r}{r+s}, s = 0, 1, \dots,$ conjecture holds

Lower bounds for matchings in regular non-bipartite graphs

Petersen's THM: A bridgeless cubic graph has a perfect match

Problem: Find the minimum of the biggest match in $\mathcal{G}(r, 2n)$ for $r > 2$.

Does every $G \in \mathcal{G}(r, 2n)$ has a match of size $\lfloor \frac{2n}{3} \rfloor$? (True for $r = 2$.)

Esperet-Kardos-King-Král-Norine:

Every cubic bridgeless graph has at least $2^{\frac{|V|}{3656}}$ perfect matchings

An analog the van der Waerden conjecture

THM Edmonds 1965: A symmetric doubly stochastic matrix with zero diagonal of even order $A = [a_{ij}]_{i,j=1}^{2n}$ is a convex combination of symmetric permutation matrices with zero diagonal if and only if $\sum_{i,j \in S} a_{ij} \leq |S| - 1$ for any odd subset $S \subset \{1, \dots, 2n\}$ (*)

Denote by Ψ_{2n} the subset of all symmetric doubly stochastic matrices of the above form

Problem: Find $\min \text{haf}(A), A \in \Psi_{2n}$

CONJECTURE: The minimum is achieved only for the matrix $\frac{1}{2n-1}A(K_{2n})$

$$\text{haf}\left(\frac{1}{2n-1}A(K_{2n})\right) \approx e^{-n}\sqrt{2e} < \text{haf}\left(\frac{1}{n}A(K_{n,n})\right) \approx e^{-n}\sqrt{2\pi n}$$

An analog the Tverberg conjecture

$$\min_{A \in \Psi_{2n}} \text{haf}_k(A) = \text{haf}_k\left(\frac{1}{2n-1}A(K_{2n})\right) =$$

$$\binom{2n}{2k} \frac{1}{(2n-1)^k} \frac{(2k)!}{2^k k!}$$

Note $\frac{1}{3}A(G) \in \Psi_{2n}$, $G \in \mathcal{G}(3, 2n)$ iff there are at least 3 edges coming out of any the set of odd number of vertices (Improves significantly the lower bound [5] if gen. v.d. Waerden conj. true)

Hyperbolic polynomials

THM: Good lower bounds hold for $\text{haf}_k(A)$ if $A \in \Psi_{2n}$ $n-1$ $n-1$ eigenvalues of A are nonpositive

Outline of proof: Fact $\mathbf{x}^\top \mathbf{A} \mathbf{x}$ is a hyperbolic polynomial for a nonnegative symmetric matrix iff A has all but one nonpositive eigenvalues [4]

$$\text{haf}_k A = (2^k k!)^{-1} \sum_{1 \leq i_1 < \dots < i_{2k} \leq 2n} \frac{\partial^{2k}}{\partial x_{i_1} \dots \partial x_{i_{2k}}} (\mathbf{x}^\top \mathbf{A} \mathbf{x})^k$$

Use the arguments of [1] to show

$$\text{haf } A \geq \left(\frac{n-1}{n}\right)^{(n-1)n} \geq e^{-n}$$

Probabilistic Methods I

$A = [a_{ij}] \in \mathbb{R}_+^{n \times n}$, $X(A) := [\sqrt{a_{ij}}x_{ij}]$,
 x_{ij} independent random variables $E(x_{ij}) = 0$, $E(x_{ij}^2) = 1$
 $E((\det X(A))^2) = \text{perm } A$. Godsil-Gutman 1981

Concentration results

A. Barvinok 1999 -

1. x_{ij} real Gaussian $\Rightarrow \det X(A)^2$ with high probability
 $\in [c^n \text{perm } A, \text{perm } A]$ $c \approx 0.28$
2. x_{ij} complex Gaussian $E(|x_{ij}|^2) = 1 \Rightarrow |\det X(A)|^2$ with high probability
 $\in [c^n \text{perm } A, \text{perm } A]$ $c \approx 0.56$
3. x_{ij} quaternion Gaussian $E(|x_{ij}|^2) = 1 \Rightarrow |\det X(A)|^2$ with high probability
 $\in [c^n \text{perm } A, \text{perm } A]$ $c \approx 0.76$

Friedland-Rider-Zeitouni 2004:

$0 < a \leq a_{ij} \leq b$, x_{ij} real Gaussian $\Rightarrow \det X(A)^2$ with high probability
 $\in [(1 - \varepsilon_n) \text{perm } A, \text{perm } A]$ $\varepsilon_n \rightarrow 0$

FRZ results use concentration for $\log_{\varepsilon} \det Z(A) = \text{tr } f(Z(A))$,
 $Z(A) = X(A)^{\top} X(A) \succeq 0$, $f = \log_{\varepsilon} x = \log \max(x, \varepsilon)$.

or $\log_{\varepsilon} \det Y(A)$, $Y(A) = \begin{bmatrix} 0 & X(A) \\ -X(A)^{\top} & 0 \end{bmatrix}$

$E(\det(\sqrt{t}I + Y(A))) = \Phi_{G_w}(t)$, $t \geq 0$

matching polynomial for weighted graph induced by A

Thm: Concentration of $\log \det(\sqrt{t}I + Y(A))$ **around expected value**

$\log \tilde{\Phi}_{G_w}(t)$, $t > 0$ **which less** $\log \Phi_{G_w}(t)$

$$\frac{1}{n} \log \tilde{\Phi}(t, G_w) \leq \frac{1}{n} \log \Phi(t, G_w) \leq \frac{1}{n} \log \tilde{\Phi}(t, G_w) + \min\left(\frac{\max_{i,j} |a_{ij}|}{2t}, 1.271\right)$$

Meaning of $\tilde{\Phi}_{G_w}(t)$?

Prob. Methods III-Matching in nonbipartite graphs

Make each undirected edge (i, j) with weight $a_{ij} = a_{ji} \geq 0$ to two opposite directed edges with weights $\pm a_{ij}$ to obtain a skew symmetric matrix

$$B = [b_{ij}] \in \mathbb{R}^{(2n) \times (2n)}, b_{ij} = 0$$

$$Y(B) = [\text{sign}(b_{ij}) \sqrt{|b_{ij}|} x_{ij}], x_{ij} = x_{ji}, x_{12}, \dots, x_{(2n-1), (2n)} \text{ i.r.v}$$

$$E(x_{ij}) = 0, E(x_{ij}^2) = 1$$

$$E(\det Y(B)) = \text{haf} A -$$

total weight of weighted matchings in induced graph by A

$$E(\det(\sqrt{t}I + Y(B))) = \Phi_{G_w}(t) - \text{the weighted matching polynomial of } G(A).$$

All the results for bipartite graphs carry over to nonbipartite graphs

Jerrum-Sinclair-Vigoda 2004: fully polynomial randomized approximation scheme (fpras) to compute $\text{perm } A$

A variation of MCMC method using rapidly mixed Markov chains converging to equilibrium point

The proofs do not carry over for nonbipartite graphs

Any $\#P$ complete problem has fpras?

Expected values of k -matchings for bipartite graphs

- **Permutation** $\sigma : \langle nr \rangle \rightarrow \langle nr \rangle$ induces $\mathbf{G}(\sigma) \in \mathcal{GB}_{\text{mult}}(r, 2n)$ and vice versa

$$\mathbf{G}(\sigma) = \left\{ \left(i, \left\lceil \frac{\sigma((i-1)r+j)}{r} \right\rceil \right), j = 1, \dots, r, i = 1, \dots, n \right\} \subset \langle n \rangle \times \langle n \rangle$$

number of different σ inducing the same simple \mathbf{G} is $(r!)^n$

- μ probability measure on $\mathcal{GB}_{\text{mult}}(r, 2n)$:

$$\mu(\mathbf{G}(\sigma)) = ((nr)!)^{-1}$$

- **FKM 06**:

$$E(k, n, r) := \mathbb{E}(\phi(k, \mathbf{G})) = \binom{n}{k}^2 r^{2k} k! (nr - k)! (nr!)^{-1},$$

$$k = 1, \dots, n$$

- $1 \leq k_l \leq n_l, l = 1, \dots$, increasing sequences of integers s.t.









$$\lim_{l \rightarrow \infty} \frac{k_l}{n_l} = p \in [0, 1]. \text{ Then}$$

$$\lim_{l \rightarrow \infty} \frac{\log E(k_l, n_l, r)}{2n_k} = f(p, r)$$









$$f(p, r) := \frac{1}{2} (p \log r - p \log p - 2(1-p) \log(1-p) + (r-p) \log(1 - \frac{p}{r}))$$

Similar results for non-bipartite graphs?







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