## Tensors

Shmuel Friedland Univ. Illinois at Chicago

University of Kansas, August 20, 2010

## Foreword

In the past ten years, tensors again became a hot topic of research in pure and applied mathematics. In applied mathematics it is driven by data which has a few parameters. In pure math. it is quantum information theory, and multilinear algebra. There are many interesting numerical and theoretical problems that need to be resolved. Tensors are related to matrices one one hand and on the other hand are related to polynomial maps.

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To paraphrase Max Noether:
Matrices were created by God and tensors by Devil.

## Overview

Ranks of 3-tensors

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(1) Basic facts.

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(2) Complexity.

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Characterization of tensor in $\mathbb{C}^{4 \times 4 \times 4}$ of border rank 4

## Basic notions

scalar $a \in \mathbb{F}$, vector $\mathbf{x}=\left(x_{1}, \ldots, x_{n}\right)^{\top} \in \mathbb{F}^{n}$, matrix $A=\left[a_{i j}\right] \in \mathbb{F}^{m \times n}$, 3-tensor $\mathcal{T}=\left[t_{i, j, k}\right] \in \mathbb{F}^{m \times n \times 1}$, p-tensor $\mathcal{T}=\left[t_{i_{1}, \ldots, i_{p}}\right] \in \mathbb{F}^{n_{1} \times \ldots \times n_{p}}$

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rank $\mathcal{T}$ minimal $r$ :
$\mathcal{T}=f_{r}\left(\mathbf{x}_{1}, \mathbf{y}_{1}, \mathbf{z}_{1}, \ldots, \mathbf{x}_{r}, \mathbf{y}_{r}, \mathbf{z}_{r}\right):=\sum_{i=1}^{r} \mathbf{x}_{i} \otimes \mathbf{y}_{i} \otimes \mathbf{z}_{i}$,

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COR $\operatorname{rank} \mathcal{T} \leq \min (m n, m l, n l)$

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PRF: 3-sat with $n$ variables $m$ clauses
satisfiable iff $\left.\operatorname{rank} \mathcal{T}=4 n+2 m, \mathcal{T} \in \mathbb{F}^{(2 n+3 m) \times(3 n) \times(3 n+m)}\right)$ otherwise rank is larger

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Border rank of $\mathcal{T}$ the minimum $k$ s.t. $\mathcal{T}$ is a limit of $\mathcal{T}_{j}, j \in \mathbb{N}$, rank $T_{j}=k$. generic rank is the rank of a random tensor $\mathcal{T} \in \mathbb{C}^{m \times n \times I}: \operatorname{grank}(m, n, l)$

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typical rank takes all the values $k=\operatorname{grank}(m, n, l), \ldots, \operatorname{mtrank}(m, n, l)$
In all the examples we know $\operatorname{mtrank}(m, n, I) \leq \operatorname{grank}(m, n, I)+1$

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Reason: For $I=(m-1)(n-1)+1$ a generic subspace of matrices of dimension / in $\mathbb{C}^{m \times n}$ intersect the variety of rank one matrices in $\mathbb{C}^{m \times n}$ at least at / lines which contain / linearly independent matrices

COR: $\operatorname{grank}_{\mathbb{C}}(2, n, I)=\min (I, 2 n)$ for $2 \leq n \leq I$
Dimension count for $\mathbb{F}=\mathbb{C}$ and $2 \leq m \leq n \leq I \leq(m-1)(n-1)+1$ :
$f_{r}:\left(\mathbb{C}^{m} \times \mathbb{C}^{n} \times \mathbb{C}^{\prime}\right)^{r} \rightarrow \mathbb{C}^{m \times n \times I}, \mathbf{x} \otimes \mathbf{y} \otimes \mathbf{z}=(a \mathbf{x}) \otimes(b \mathbf{y}) \otimes\left((a b)^{-1} \mathbf{z}\right)$

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Product of two $2 \times 2$ matrices is done by 7 multiplications

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$\operatorname{grank}(3,2 p, 2 p)=\left\lceil\frac{12 p^{2}}{4 p+1}\right\rceil$ and $\operatorname{grank}(3,2 p-1,2 p-1)=\left\lceil\frac{3(2 p-1)^{2}}{4 p-1}\right\rceil+1$

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Avoid round-off error:
$\mathbf{w}_{r} \in\left(\mathbb{Z}^{m} \times \mathbb{Z}^{n} \times \mathbb{Z}^{\prime}\right)^{r}$ find rank $J\left(f_{r}\right)\left(\mathbf{w}_{r}\right)$ exact arithmetic

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I checked the conjecture up to $m, n, I \leq 14$

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For $m n \leq I \operatorname{mtrank}(m, n, I)=\operatorname{grank}(m, n, I)=m n$.
For $2 \leq m \leq n \leq I<m n-1$, there exist $V_{1}, \ldots, V_{c(m, n, l)} \subset \mathbb{R}^{m \times n \times I}$ pairwise distinct open connected semi-algebraic sets s.t.

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$\mathbb{R}^{m \times n \times I} \operatorname{IPS}:\langle\mathcal{A}, \mathcal{B}\rangle=\sum_{i=j=k}^{m, n, l} a_{i, j, k} b_{i, j, k},\|\mathcal{T}\|=\sqrt{\langle\mathcal{T}, \mathcal{T}\rangle}$

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\begin{aligned}
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& \langle\mathbf{x} \otimes \mathbf{y} \otimes \mathbf{z}, \mathbf{u} \otimes \mathbf{v} \otimes \mathbf{w}\rangle=\left(\mathbf{u}^{\top} \mathbf{x}\right)\left(\mathbf{v}^{\top} \mathbf{y}\right)\left(\mathbf{w}^{\top} \mathbf{z}\right)
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Equivalent: $\max _{\|\mathbf{x}\|=\|\mathbf{y}\|=\|\mathbf{z}\|=1} \sum_{i=j=k}^{m, n, l} t_{i, j, k} x_{i} y_{j} z_{k}$

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$\lambda$ singular value, $\mathbf{x}, \mathbf{y}, \mathbf{z}$ singular vectors

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$\lambda$ singular value, $\mathbf{x}, \mathbf{y}, \mathbf{z}$ singular vectors
How many distinct singular values are for a generic tensor?

## $\ell_{p}$ maximal problem and Perron-Frobenius

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$$
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$$

## $\ell_{p}$ maximal problem and Perron-Frobenius

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$p=3$ is most natural in view of homogeneity
Assume that $\mathcal{T} \geq 0$. Then $\mathbf{x}, \mathbf{y}, \mathbf{z} \geq 0$

For which values of $p$ we have an analog of Perron-Frobenius theorem?

Yes, for $p \geq 3$, No, for $p<3$,
Friedland-Gauber-Han [1]

## $\left(R_{1}, R_{2}, R_{3}\right)$-rank approximation of 3-tensors

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Fundamental problem in applications:
Approximate well and fast $\mathcal{T} \in \mathbb{R}^{m_{1} \times m_{2} \times m_{3}}$ by rank $\left(R_{1}, R_{2}, R_{3}\right)$ 3 -tensor.

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## $\left(R_{1}, R_{2}, R_{3}\right)$-rank approximation of 3-tensors

Fundamental problem in applications:
Approximate well and fast $\mathcal{T} \in \mathbb{R}^{m_{1} \times m_{2} \times m_{3}}$ by rank $\left(R_{1}, R_{2}, R_{3}\right)$ 3 -tensor.

Best ( $R_{1}, R_{2}, R_{3}$ ) approximation problem:
Find $\mathbb{U}_{i} \subset \mathbb{F}^{m_{i}}$ of dimension $R_{i}$ for $i=1,2,3$ with maximal $\left\|P_{\mathbb{U}_{1} \otimes \mathrm{U}_{2} \otimes \mathrm{U}_{3}}(\mathcal{T})\right\|$.

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Use Newton method on Grassmannians - Eldén-Savas 2009 [1]

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CUR approximation of $\mathcal{A}$ obtained by choosing $E, F, G$ submatrices of unfolded $\mathcal{A}$ in the mode $1,2,3$.

## List of applications

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Factor analysis

## Scaling of nonnegative tensors to tensors with given row, column and depth sums

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& 0 \leq \mathcal{T}=\left[t_{i, j, k}\right] \in \mathbb{R}^{m \times n \times I} \text { has given row, column and depth sums: } \\
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$f_{\mathcal{T}}$ is convex
$f_{\mathcal{T}}$ is strictly convex implies $\mathcal{T}$ is not decomposable: $\mathcal{T} \neq \mathcal{T}_{1} \oplus_{\mathcal{I}} \mathcal{T}_{2}$.

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## Characterization of tensor in $\mathbb{C}^{4 \times 4 \times 4}$ of border rank 4

Major problem in algebraic statistics: phylogenic trees and their invariants [1]

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Strassen's condition hold for any $3 \times 3 \times 3$ subtensor of $\mathcal{T}$ : $\operatorname{det}(U(\operatorname{adj} W) V-V(\operatorname{adj} W) U)=0, \quad U, V, W \in \mathbb{C}^{3 \times 3 \times 3}$ equations of degree 9

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If $\mathbf{W}$ contains an invertible matrix $Z$ then any other $X, Y \in \mathbf{W}$ satisfy $X(\operatorname{adj} Z) Y=Y(\operatorname{adj} Z) X$ - equations of degree 5
 equations of degree 6

Strassen's condition hold for any $3 \times 3 \times 3$ subtensor of $\mathcal{T}$ : $\operatorname{det}(U(\operatorname{adj} W) V-V(\operatorname{adj} W) U)=0, \quad U, V, W \in \mathbb{C}^{3 \times 3 \times 3}$ equations of degree 9

Friedland [5] one needs a equations of degree 16

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