

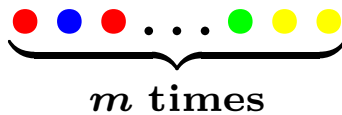
**MULTI-DIMENSIONAL ENTROPY
and the monomer-dimer problem**

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1 UNRESTRICTED CAPACITY

$$\langle n \rangle := \{1, 2, 3, \dots, n\}$$

IS AN ALPHABET ON n LETTERS.



LINEAR STORAGE OF LENGTH m .

STORAGE MESSAGES n^m .

(UNRESTRICTED) CAPACITY:

$$\log_2 n^m = \frac{\log_2 n^m}{m}.$$

2 (0 - 1) LIMITED CHANNEL

(HARD CORE LATTICE or NEAR NEIGHBOR EXCLUSION)

$$n = 2, \langle 2 \rangle = \{1, 2\} = \{1, 0\} \quad (2 \equiv 0).$$

NO TWO 1's ARE NEIGHBORS.

OF m -MESSAGES IS u_m ;

$$u_{m+1} = u_m + u_{m-1}, \quad m = 1, 2, \dots$$

FIBONACCI SEQUENCE 2, 3, 5, 8, ...

SHANNON CAPACITY:

$$\lim_{m \rightarrow \infty} \frac{\log_2 u_m}{m} =$$
$$\log_2 \frac{1 + \sqrt{5}}{2} = 0.694241914\dots$$

3 SUBSHIFT OF FINITE TYPE

$$a = a_1 \dots a_m = (a_i)_1^m : \langle m \rangle \rightarrow \langle n \rangle$$

WORD OF LENGTH m .

$\langle n \rangle^{\langle m \rangle}$ all messages of length m .

$$\langle n \rangle^{\mathbb{Z}} = \{a = (\dots a_{-1} a_0 a_1 \dots)\} =$$

$$\{(a_i)_{i \in \mathbb{Z}} : \mathbb{Z} \rightarrow \langle n \rangle\}$$

$$\langle n \rangle^{\mathbb{N}} = \{a = (a_1 \dots)\} =$$

$$\{(a_i)_{i \in \mathbb{N}} : \mathbb{N} \rightarrow \langle n \rangle\}$$

PROJECTION: $\pi_m((a_i)) = (a_i)_{i=1}^m$.

$\mathcal{S} \subset \langle n \rangle^{\mathbb{N}}$ ($\langle n \rangle^{\mathbb{Z}}$) is SOFT if

$\exists P \subset \langle n \rangle^{\langle r \rangle}$ such that $a \in \mathcal{S} \iff$

any consec. string of r letters in a in P .

EXAMPLE: $r = 2 \Rightarrow P = \Gamma \subset \langle n \rangle \times \langle n \rangle$.

$$n = 2, P = \{\bullet\bullet, \bullet\bullet, \bullet\bullet\}$$



ALLOWABLE WORD OF LENGTH $m + 1$

A WALK OF LENGTH m ON Γ

$$\Gamma^m = \{(a_i)_1^{m+1} \in \langle n \rangle^{m+1} :$$

$$(a_i, a_{i+1}) \in \Gamma\}$$

$$\Gamma^{\mathbb{N}} = \{(a_i) \in \langle n \rangle^{\mathbb{N}} : (a_i, a_{i+1}) \in \Gamma\}$$

$$\Gamma^{\mathbb{Z}} = \{(a_i) \in \langle n \rangle^{\mathbb{Z}} : (a_i, a_{i+1}) \in \Gamma\}$$

any **SOFT** can be coded as a walk on

$$\Gamma \subset \langle N \rangle \times \langle N \rangle \equiv \pi_{r-1}(P) \times \pi_{r-1}(P)$$

$$N = \#\pi_{r-1}(P)$$

$$((a_i)_1^{r-1}, (b_i)_1^{r-1}) \in \Gamma \iff$$

$$b_1 = a_2, \dots, b_{r-2} = a_{r-1},$$

ASSUMPTION: \mathcal{S} -SOFT

$(a_i)_{i \in \mathbb{N}}$ is m -periodic

$$a_{i+m} = a_i, \quad i \in \mathbb{N}.$$

$\delta_m = \log \# P$ -allow. words of length m .

$$\tilde{\delta}_m = \log \# \pi_m(\mathcal{S})$$

$\delta_{m,per} = \log \# m$ -periodic words in \mathcal{S} .

$$\delta_{m,per} \leq \tilde{\delta}_m \leq \delta_m.$$

$\{t_i\}_{i \in \mathbb{N}} \subset \mathbb{R}$ IS SUBADDITIVE (SA):

$$t_{p+q} \leq t_p + t_q \quad \text{for all } p, q \in \mathbb{N}.$$

$$\{t_i\}_{i \in \mathbb{N}} \text{ SA} \Rightarrow \lim_{i \rightarrow \infty} \frac{t_i}{i} = \tau \leq \frac{t_p}{p}.$$

CLAIM: $\{\delta_m\}, \{\tilde{\delta}_m\}$ - ARE SA.

$$h_{com} := \lim_{m \rightarrow \infty} \frac{\delta_m}{m} \quad \text{CAPACITY}$$

$$h := \lim_{m \rightarrow \infty} \frac{\tilde{\delta}_m}{m} \quad \text{ENTROPY}$$

$$h_{per} := \limsup_{m \rightarrow \infty} \frac{\delta_{m,per}}{m} \quad \text{periodic entropy}$$

$$-\infty \leq h_{per} \leq h \leq h_{com}$$

MAIN THM

for 1- dimensional SOFT \mathcal{S}

$$h_{per} = h = h_{com} = \log \rho(\Gamma)$$

Γ -GRAPH INDUCED BY SOFT.

COR. \mathcal{S} is decidable:

EITHER $\mathcal{S} = \emptyset \iff$:

$\exists m$ with no allow. word of length m

OR \mathcal{S} contains an m periodic state.

4 MULTI-DIMENSIONAL CAPACITY

$2 \leq d$ dimension

$$\mathbf{m} := (m_1, \dots, m_d) \in \mathbb{Z}^d$$

$$|\mathbf{m}| = |m_1| + \dots + |m_d|$$

$$|\mathbf{m}|_{pr} := |m_1| \times \dots \times |m_d|$$

$$\langle \mathbf{m} \rangle := \langle m_1 \rangle \times \dots \times \langle m_d \rangle,$$

for $\mathbf{m} \in \mathbb{N}^d$

$a : \langle \mathbf{m} \rangle \rightarrow \langle \mathbf{n} \rangle$ is $(a_i)_{i \in \langle \mathbf{m} \rangle}$

$\langle \mathbf{n} \rangle^{\langle \mathbf{m} \rangle}$ set of all maps a .

$\langle \mathbf{n} \rangle^{\mathbb{N}^d}$ & $\langle \mathbf{n} \rangle^{\mathbb{Z}^d}$ all maps

from \mathbb{N}^d & \mathbb{Z}^d to $\langle \mathbf{n} \rangle$.

$\pi_{\mathbf{m}}((a_i)) = (a_i)_{i \in \langle \mathbf{m} \rangle}$ proj. on $\langle \mathbf{m} \rangle$.

$\mathcal{S} \subset \langle \mathbf{n} \rangle^{\mathbb{N}^d}$ ($\langle \mathbf{n} \rangle^{\mathbb{Z}^d}$) SOFT if

$\exists P \subset \langle \mathbf{n} \rangle^{\langle \mathbf{r} \rangle}$ such that $a \in \mathcal{S} \iff$

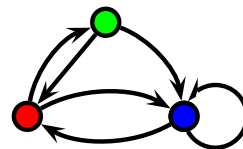
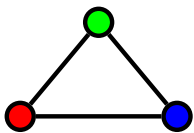
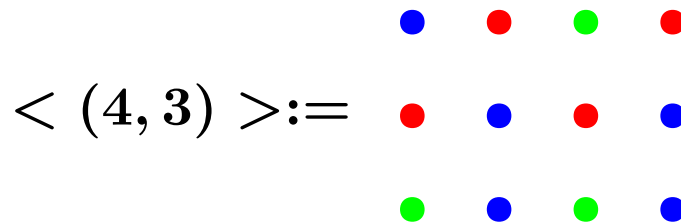
any consec. box of \mathbf{r} letters in a in P .

EXAMPLE: $\Gamma = (\Gamma_1, \dots, \Gamma_d)$

$(a_i)_{i \in \langle m \rangle}$ IS ALLOWABLE WORD

if any line through $i \in \langle m \rangle$

in direction $e_k = (\delta_{k1}, \dots, \delta_{kd})$ is in $\Gamma_k^{m_k-1}$ for each k



any **SOFT** can be coded as $\Gamma^{\mathbb{N}^d} (\Gamma^{\mathbb{Z}^d})$

$\exists \mathcal{S}$ **SOFT NOT DECIDABLE:**

$\exists \Gamma^{\mathbb{N}^2} \neq \emptyset$ with no periodic state

Berger 1966

$$\Gamma^{\mathbb{N}^d} = \emptyset \iff \exists m \Gamma^m = \emptyset$$

ASSUMPTION: \mathcal{S} -SOFT

δ_m -log# P -allow. words of dim. m .

$$\tilde{\delta}_m = \log \# \pi_m(\mathcal{S})$$

$\delta_{m,per}$ -log# m -periodic words in \mathcal{S} .

$$\delta_{m,per} \leq \tilde{\delta}_m \leq \delta_m.$$

$\{\delta_m\}, \{\tilde{\delta}_m\}$ - are SA in each coordinate

(split box $\langle m \rangle$ to 2 boxes by $x_k = i_k$)

$$h_{com} := \lim_{m \rightarrow \infty} \frac{\delta_m}{|m|_{pr}} \quad \text{CAPACITY}$$

$$h := \lim_{m \rightarrow \infty} \frac{\tilde{\delta}_m}{|m|_{pr}} \quad \text{ENTROPY}$$

$$h_{per} := \limsup_{m \rightarrow \infty} \frac{\delta_{m,per}}{|m|_{per}} \quad \text{PERIODIC ENTROPY}$$

$$-\infty \leq h_{per} \leq h \leq h_{com} \leq \frac{\delta_m}{|m|_{pr}}$$

Berger: $-\infty = h_{per} < 0 \leq h \leq h_{com}$

Friedland 97: for all \mathcal{S} $h = h_{com}$

5 UPPER ESTIMATES OF MDC

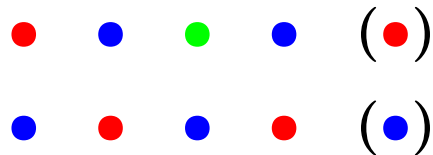
$d = 2$ and $\mathbf{m} = (m_1, m_2)$.

$$\Gamma(2, m_1) \subset \Gamma_1^{m_1-1} \times \Gamma_1^{m_1-1}$$

$$((b_i)_1^{m_1}, (c_i)_1^{m_1}) \in \Gamma(2, m_1) \iff (b_i, c_i) \in \Gamma_2$$

for each $i = 1, \dots, m_1$.

EXAMPLE: $\Gamma(2, 4)$



$\Gamma_{per}(2, m_1) \subset \Gamma(2, m_1)$:

$$(b_{m_1}, b_1), (c_{m_1}, c_1) \in \Gamma_1$$

$$\rho(\Gamma(2, m_1)) \geq \rho(\Gamma_{per}(2, m_1))$$

statistical mechanics - *transfer* matrices:

$A(\Gamma(\mathbf{2}, m_1))$ & $A(\Gamma_{per}(\mathbf{2}, m_1))$

$$\lim_{m_2 \rightarrow \infty} \frac{\delta_m}{m_2} = \log \rho(\Gamma(\mathbf{2}, m_1)) \geq m_1 h$$

$$\limsup_{m_2 \rightarrow \infty} \frac{\delta_{m,per}}{m_2} = \log \rho(\Gamma_{per}(\mathbf{2}, m_1))$$

$$\lim_{m_1 \rightarrow \infty} \frac{\log \rho(\Gamma(\mathbf{2}, m_1))}{m_1} = h$$

$$\limsup_{m_1 \rightarrow \infty} \frac{\log \rho(\Gamma_{per}(\mathbf{2}, m_1))}{m_1} = h_{per}$$

6 Upper-Lower Bounds with Symmetricity Con.

Γ_2 - SYMMETRIC

$$\max\left(\frac{\log \rho(\Gamma(1, p + 2q + 1)) - \log \rho(\Gamma(1, 2q + 1))}{p}, \right.$$

$$\left. \frac{\log \rho(\Gamma_{per}(1, p + 2q)) - \log \rho(\Gamma_{per}(1, 2q))}{p} \right) \leq$$

$$h \leq \frac{\log \rho(\Gamma_{per}(1, 2m))}{2m} \left(\leq \frac{\log \rho(\Gamma(1, 2m))}{2m} \right)$$

for any $m, p \geq 1$ and $q \geq 0$.

\mathcal{S} DECIDABLE:

$$\mathcal{S} \neq \emptyset \iff \Gamma(1, 2) \text{ has cycle}$$

AND COMPUTABLE

$$h_{per} = h$$

$d \geq 3$ $\Gamma_1, \dots, \Gamma_{d-1}$ symmetric \Rightarrow

$h_{per} = h$ is computable

$$\frac{\log \rho(\Gamma_{per}((2m_1, \dots, 2m_{d-1})))}{2^{d-1} |m_1| \dots |m_{d-1}|} \geq h \geq$$

$$\frac{\bar{h}(p + 2q) - \bar{h}(2q)}{p}$$

$$\bar{h}(q) := \lim_{m^{\{1\}} \rightarrow \infty} \log \frac{\#W_{\{1\}, per}(m^{\{1\}}, q)}{|m^{\{1\}}|_{pr}}$$

TRUE IF Γ ISOTROPIC & SYMMETRIC:

$$\Gamma_1 = \dots = \Gamma_d = \Delta - \text{SYMMETRIC}$$

7 Automorphism Subgroups and Computations

$A = (a_{ij})_1^N$ nonnegative matrix

$\mathcal{A}(A) := \{\pi \in S_N : a_{\pi(i)\pi(j)} = a_{ij}, i, j \in \langle N \rangle\}$

$G \leq \mathcal{A}(A), \mathcal{O}(G) := \langle N \rangle / G, M = \#\mathcal{O}(G)$

$$\hat{A} = (\hat{a}_{\alpha\beta})_{\alpha, \beta \in \mathcal{O}(G)},$$

$$\hat{a}_{\alpha\beta} =: \sum_{j \in \beta} a_{ij}, i \in \alpha,$$

$$\rho(A) = \rho(\hat{A}),$$

$A = A^T \Rightarrow \hat{A}$ is symmetric for

$$\langle x, y \rangle = \sum_{\alpha \in \mathcal{O}(G)} (\#\alpha) x_\alpha y_\alpha.$$

$$M \geq N/\#G,$$

In our computations $M \sim N/\#G$.

$$T((m_1, \dots, m_d)) := (\mathbb{Z}/m_1\mathbb{Z}) \times \dots \times (\mathbb{Z}/m_d\mathbb{Z})$$

$$\mathcal{A}(m) = \mathcal{A}(\text{adjacency graph of } T(m))$$

Γ is isotropic symmetric graph \Rightarrow

$$\mathcal{A}(m^{\{d\}}) \leq \mathcal{A}(\mathcal{A}(\Gamma_{per}(d, m^{\{d\}})))$$

$$\#\mathcal{A}(2) = 2, \#\mathcal{A}(m) = 2m, m \geq 3 \Rightarrow$$

$$\mathcal{A}(m_1) \times \dots \times \mathcal{A}(m_{d-1}) \leq \mathcal{A}(m^d)$$

$$m_i = m_j \Rightarrow i \leftrightarrow j \Rightarrow$$

$$\#\mathcal{A}((m, \dots, m)) \geq (2m)^{d-1} (d-1)! \text{ for } m \geq 3$$

$$T(4) \sim \langle (2, 2) \rangle \Rightarrow$$

$$\#\mathcal{A}(4, \dots, 4) \geq 2^{2(d-1)} (2(d-1))!$$

Friedland-Peled 03

8 d -Dimensional Monomer-Dimers

Dimer: (\mathbf{i}, \mathbf{j}) , $\mathbf{j} = \mathbf{i} + \mathbf{e}_k \in \mathbb{Z}^d$.

any partition of \mathbb{Z}^d to dimers (1-factor).

Monomer: occupies $\mathbf{i} \in \mathbb{Z}^d$.

any partition of \mathbb{Z}^d to monomer-dimers

is 1-factor of a subset of \mathbb{Z}^d .

Dimer and Monomer-Dimer are SOFT

$$0 = \tilde{h}_1 \leq \tilde{h}_2 \leq \dots \leq \tilde{h}_d \leq \dots (\text{dimers})$$

$$\log \frac{1 + \sqrt{5}}{2} = h_1 \leq h_2 \leq \dots \leq h_d \leq \dots$$

(monomer – dimer)

Fisher, Kasteleyn and Temperley 61

$$\tilde{h}_2 = \frac{1}{\pi} \sum_{i=0}^{\infty} \frac{(-1)^i}{(2i+1)^2} = 0.29156090\dots$$

9 Upper and Lower Bounds

$A = (a_{ST}), S, T \subset T(\mathbf{m}^{\{d\}})$ is transfer matrix
 represents multigraph and is "Hankel": a_{ST} is
 the number of monomer-dimer (dimer) covers of
 $T(\mathbf{m}^{\{d\}}) \setminus (S \cup T)$ and $a_{ST} = 0$ if $S \cap T \neq \emptyset$
 S (T) the locations of dimers going down (up)

$$\tilde{h}_d \leq \frac{\log \tilde{\beta}(2\mathbf{m}^{\{d\}})}{2^{d-1} |\mathbf{m}^{\{d\}}|_{pr}} \quad \text{Ciucu 98}$$

$$h_d \leq \frac{\log \beta(2\mathbf{m}^{\{d\}})}{2^{d-1} |\mathbf{m}^{\{d\}}|_{pr}} \quad \text{FP 03}$$

$$h_2 \geq \frac{\log \beta(p + 2q) - \log \beta(2q)}{p} \quad \text{FP 03}$$

Similar lower bounds for exist for h_3, \tilde{h}_3 FP 03.

$$\tilde{h}_3 \leq \frac{\log \tilde{\beta}(6,4)}{6 \cdot 4} = 0.4575469308$$

Lundow 2001 using matrix of order 184854

which splits to a direct sum of 3 matrices.

10 The Value of h_2

Friedland-Peled 03

(confirming Baxter's heuristic computations 1968):

$$h_2 = 0.66279897 \text{ using } m = 14, 15, 16$$

$$.66279897190 \leq h_2 \leq .662798972844913$$

$$m = 16, N = 2^{16} = 65536,$$

$$M = 2250, \frac{N}{2m} = 2,048.$$

M. Jerrum-87:

Two-dimensional monomer-dimer systems
are computationally intractable.

m_1	$\#\mathcal{O}(m_1)$	$\log \beta(m_1)$
4	6	2.6532941163
5	8	3.3135066910
6	13	3.9769139475
7	18	4.6395628723
8	30	5.3023993987
9	46	5.9651887945
10	78	6.6279902386
11	126	7.2907885674
12	224	7.9535877093
13	380	8.6163866375
14	687	9.2791856222
15	1224	9.9419845918
16	2250	10.60478356551861
17	4112	$\in (11.267582535, 11.267582554)$

Table 1: Spectral radii for h_2

11 Matchings and permanents

$\text{perm}_s A = \text{sum perm. all } s \times s \text{ submatrices of } A.$

Tverberg's conjecture (Friedland 1982)

$\text{perm}_s A \geq \text{perm}_s J_n$ for any $n \times n$ d.s. matrix.

G is r -regular bipartite with n in each class.

$W(G, s)$ set of all matchings of size s in G .

$$\#W(G, s) \geq \binom{n}{s}^2 s! \left(\frac{r}{n}\right)^s$$

$\lambda_d(p)$ -monomer dimer entropy
with dimer density $p \in [0, 1]$.

$$\lambda_d(p) \geq \frac{1}{2}(-p \log p - 2(1-p) \log(1-p) + p \log 2d - p)$$

$$h_d = \max_{p \in [0,1]} \lambda_d(p) \geq \frac{1}{2}(-p(d) \log p(d) - 2(1 - p(d)) \log(1 - p(d)) + p(d) \log 2d - p(d))$$

$$p(d) = \frac{4d+1-\sqrt{8d+1}}{4d}.$$

12 The Graphs for $\lambda_2(p)$

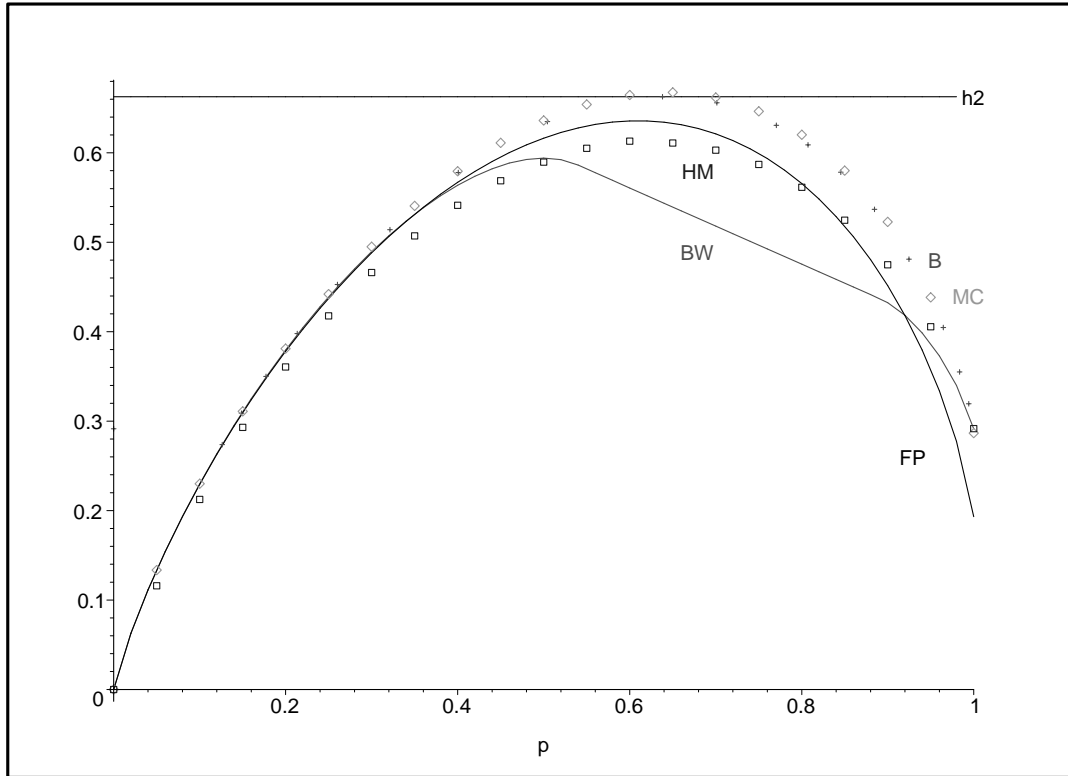


Figure 1: HM is the lower bound of Hammersley-Menon, BW is the lower bound of Bondy-Welsh, FP is the lower bound of Friedland-Peled, MC is the Monte Carlo estimate of Hammersley-Menon, B are Baxter's estimates, and h_2 is the true value of $h_2 = \max \lambda_2(p)$.

13 Estimates for \tilde{h}_3 and h_3

Schrijver 1998

$$\#W(G, n) \geq \left(\frac{(r-1)^{r-1}}{r^{r-2}} \right)^n$$

$$\tilde{h}_d = \lambda_d(1) \geq$$

$$\frac{1}{2}((2d-1) \log(2d-1) - (2d-2) \log 2d)$$

$$(S) \ 0.440075842 \leq \tilde{h}_3 \leq 0.4575469308 \ (L)$$

FP gives

$$0.7652789557 \leq h_3 \leq .7862023450$$

(m_1, m_2)	$\#\mathcal{O}(m_1, m_2)$	$\log \beta(m_1, m_2)$	$\frac{\log \beta(m_1, m_2)}{m_1 m_2}$
(2, 2)	6	3.224405658	0.8061014
(3, 2)	13	4.768958913	0.7948264
(4, 2)	34	6.367778959	0.7959723
(5, 2)	78	7.958105292	0.7958105
(6, 2)	237	9.550024542	0.7958353
(7, 2)	687	11.14163679	0.7958311
(8, 2)	2299	12.73331093	0.7958319
(3, 3)	25	7.057039652	0.7841155
(4, 3)	158	9.421594940	0.7851329
(5, 3)	708	11.77517604	0.7850117
(4, 4)	805	12.57923752	0.7862023

Table 2: Spectral radii for h_3

14 Lower Bounds for $\lambda_3(p)$

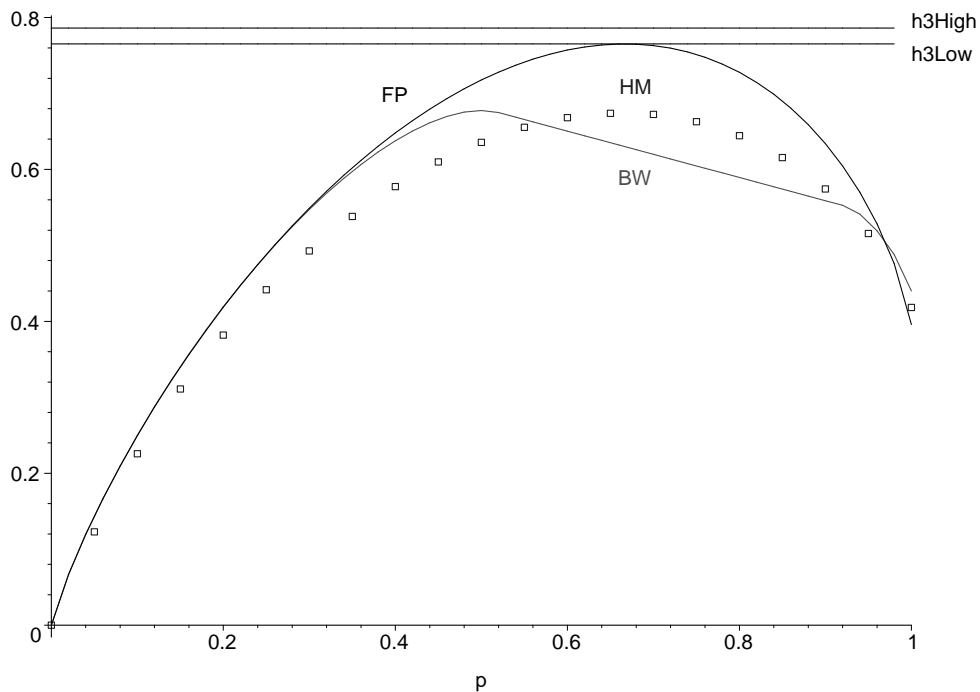


Figure 2: HM is the lower bound of Hammersley-Menon, BW is the lower bound of Bondy-Welsh, FP is the lower bound Friedland-Peled, h_{3Low} and h_{3High} are the best lower and upper bounds for $h_3 = \max \lambda_3(p)$.

15 Low. Bds \tilde{h}_3, h_3 by spec. rad.

$$h_3 \geq \frac{\log \beta(p + 2q, u + 2s) - \log \beta(p + 2q, 2s)}{up} - \frac{\log \beta(2q, 2v)}{2vp}$$

$$\tilde{h}_3 \geq \frac{\log \tilde{\beta}(p + 2q, u + 2s) - \log \tilde{\beta}(p + 2q, 2s)}{up} - \frac{\log \tilde{\beta}(2q, 2v)}{2vp}$$

$$\beta(n, 0) = \beta(0, n) = \tilde{\beta}(n, 0) = \tilde{\beta}(0, n) = 2^n$$

$$n \in \mathbb{N}$$