

Exercise set #6, exercises K and L— Summer 2007

K. The dynamic programming algorithm for finding the cost of computing $M_1 \times M_2 \times \dots \times M_{18}$ using optimal order gave the following output:

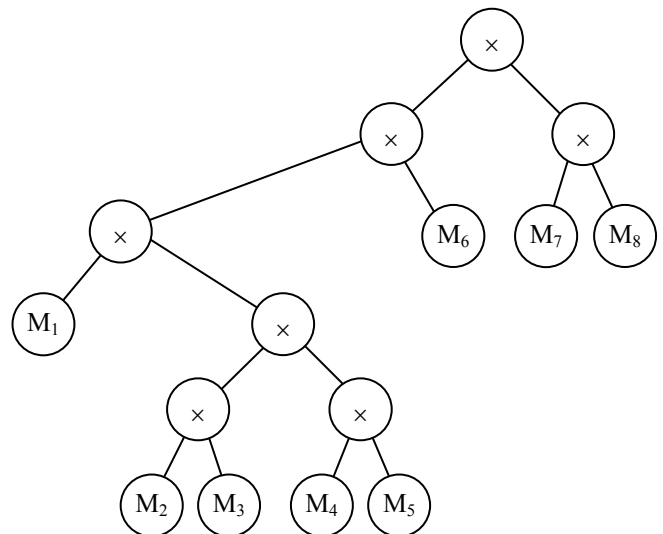
| i | j | m_{ij} | v_{ij} |
|-----|-----|----------|----------|
| 1 | 2 | 150 | 1 |
| 2 | 3 | 360 | 2 |
| 3 | 4 | 180 | 3 |
| 4 | 5 | 240 | 4 |
| 5 | 6 | 120 | 5 |
| 6 | 7 | 216 | 6 |
| 7 | 8 | 108 | 7 |
| 8 | 9 | 72 | 8 |
| 9 | 10 | 80 | 9 |
| 10 | 11 | 600 | 10 |
| 11 | 12 | 1200 | 11 |
| 12 | 13 | 360 | 12 |
| 13 | 14 | 168 | 13 |
| 14 | 15 | 231 | 14 |
| 15 | 16 | 3465 | 15 |
| 16 | 17 | 990 | 16 |
| 17 | 18 | 1080 | 17 |
| 1 | 3 | 330 | 2 |
| 2 | 4 | 330 | 2 |
| 3 | 5 | 240 | 4 |
| 4 | 6 | 480 | 4 |
| 5 | 7 | 390 | 6 |
| 6 | 8 | 156 | 6 |
| 7 | 9 | 156 | 8 |
| 8 | 10 | 260 | 8 |
| 9 | 11 | 380 | 10 |
| 10 | 12 | 1080 | 11 |
| 11 | 13 | 810 | 11 |
| 12 | 14 | 675 | 13 |
| 13 | 15 | 495 | 13 |
| 14 | 16 | 1716 | 15 |
| 15 | 17 | 1144 | 15 |
| 16 | 18 | 1254 | 17 |
| 1 | 4 | 405 | 2 |
| 2 | 5 | 360 | 2 |
| 3 | 6 | 312 | 5 |
| 4 | 7 | 888 | 5 |
| 5 | 8 | 196 | 5 |
| 6 | 9 | 188 | 8 |

| i | j | m_{ij} | v_{ij} |
|-----|-----|----------|----------|
| 7 | 10 | 308 | 8 |
| 8 | 11 | 650 | 8 |
| 9 | 12 | 620 | 11 |
| 10 | 13 | 930 | 10 |
| 11 | 14 | 1020 | 13 |
| 12 | 15 | 1086 | 13 |
| 13 | 16 | 2796 | 13 |
| 14 | 17 | 1186 | 14 |
| 15 | 18 | 1312 | 17 |
| 1 | 5 | 450 | 2 |
| 2 | 6 | 492 | 2 |
| 3 | 7 | 474 | 6 |
| 4 | 8 | 316 | 4 |
| 5 | 9 | 236 | 8 |
| 6 | 10 | 316 | 8 |
| 7 | 11 | 668 | 8 |
| 8 | 12 | 764 | 8 |
| 9 | 13 | 668 | 12 |
| 10 | 14 | 1014 | 13 |
| 11 | 15 | 1371 | 13 |
| 12 | 16 | 4101 | 13 |
| 13 | 17 | 1234 | 13 |
| 14 | 18 | 1258 | 17 |
| 1 | 6 | 552 | 2 |
| 2 | 7 | 744 | 2 |
| 3 | 8 | 388 | 3 |
| 4 | 9 | 412 | 8 |
| 5 | 10 | 376 | 8 |
| 6 | 11 | 656 | 8 |
| 7 | 12 | 824 | 8 |
| 8 | 13 | 722 | 8 |
| 9 | 14 | 710 | 13 |
| 10 | 15 | 1293 | 13 |
| 11 | 16 | 3876 | 13 |
| 12 | 17 | 1474 | 12 |
| 13 | 18 | 1426 | 17 |
| 1 | 7 | 759 | 2 |
| 2 | 8 | 448 | 2 |
| 3 | 9 | 412 | 8 |

| i | j | m_{ij} | v_{ij} |
|-----|-----|----------|----------|
| 4 | 10 | 636 | 8 |
| 5 | 11 | 726 | 8 |
| 6 | 12 | 840 | 8 |
| 7 | 13 | 812 | 8 |
| 8 | 14 | 836 | 8 |
| 9 | 15 | 864 | 14 |
| 10 | 16 | 3186 | 13 |
| 11 | 17 | 1774 | 11 |
| 12 | 18 | 1834 | 17 |
| 1 | 8 | 548 | 1 |
| 2 | 9 | 528 | 8 |
| 3 | 10 | 528 | 8 |
| 4 | 11 | 1056 | 8 |
| 5 | 12 | 896 | 8 |
| 6 | 13 | 848 | 8 |
| 7 | 14 | 902 | 8 |
| 8 | 15 | 1062 | 8 |
| 9 | 16 | 1854 | 15 |
| 10 | 17 | 1854 | 10 |
| 11 | 18 | 2014 | 17 |
| 1 | 9 | 588 | 8 |
| 2 | 10 | 728 | 8 |
| 3 | 11 | 858 | 8 |
| 4 | 12 | 1128 | 8 |
| 5 | 13 | 894 | 8 |
| 6 | 14 | 922 | 8 |
| 7 | 15 | 1104 | 8 |
| 8 | 16 | 2664 | 8 |
| 9 | 17 | 1866 | 13 |
| 10 | 18 | 1950 | 17 |
| 1 | 10 | 728 | 8 |
| 2 | 11 | 1128 | 8 |
| 3 | 12 | 1056 | 8 |
| 4 | 13 | 1056 | 8 |
| 5 | 14 | 976 | 8 |
| 6 | 15 | 1108 | 8 |
| 7 | 16 | 2502 | 8 |
| 8 | 17 | 1902 | 8 |
| 9 | 18 | 1914 | 17 |

| i | j | m_{ij} | v_{ij} |
|-----|-----|----------|----------|
| 1 | 11 | 1078 | 8 |
| 2 | 12 | 1228 | 8 |
| 3 | 13 | 1074 | 8 |
| 4 | 14 | 1194 | 8 |
| 5 | 15 | 1170 | 8 |
| 6 | 16 | 2370 | 8 |
| 7 | 17 | 1998 | 8 |
| 8 | 18 | 2118 | 17 |
| 1 | 12 | 1248 | 8 |
| 2 | 13 | 1164 | 2 |
| 3 | 14 | 1137 | 13 |
| 4 | 15 | 1444 | 8 |
| 5 | 16 | 2500 | 8 |
| 6 | 17 | 2038 | 8 |
| 7 | 18 | 2142 | 17 |
| 1 | 13 | 1246 | 8 |
| 2 | 14 | 1298 | 8 |
| 3 | 15 | 1318 | 8 |
| 4 | 16 | 3250 | 8 |
| 5 | 17 | 2078 | 5 |
| 6 | 18 | 2134 | 17 |
| 1 | 14 | 1328 | 8 |
| 2 | 15 | 1532 | 8 |
| 3 | 16 | 2512 | 8 |
| 4 | 17 | 2198 | 4 |
| 5 | 18 | 2198 | 17 |
| 1 | 15 | 1522 | 8 |
| 2 | 16 | 3202 | 8 |
| 3 | 17 | 2266 | 8 |
| 4 | 18 | 2486 | 17 |
| 1 | 16 | 2852 | 8 |
| 2 | 17 | 2326 | 2 |
| 3 | 18 | 2338 | 17 |
| 1 | 17 | 2426 | 1 |
| 2 | 18 | 2566 | 17 |
| 1 | 18 | 2546 | 17 |

Find the optimal order for multiplying $M_1 \times M_2 \times \dots \times M_{18}$. Rather than using parentheses to show the optimal order, draw a parse tree. For example, the parse tree corresponding to $((M_1 \times ((M_2 \times M_3) \times (M_4 \times M_5))) \times M_6) \times (M_7 \times M_8)$ would be as shown at right.



L. Consider the weighted digraph at right. Edges not shown have weight ∞ . In the terminology of the all-pairs shortest path algorithm for digraphs (See handout on web site.), find $d_{1,10}^k$ for $k = 0, 1, 2, \dots, 9, 10$, and find $short_k(1,10)$ for all values of k for which $d_{1,10}^k < \infty$.

Recall that $short_k(i,j)$ is a path from i to j that has minimal length, subject to the constraint that any intermediate vertices on the path lie in the set $\{1,2,\dots,k\}$, and that d_{ij}^k is the length of $short_k(i,j)$.

