## Depth-first Search in Digraphs - examples

$\Longrightarrow$ denotes a tree edge.
$d(v), f(v)$ appears above or below each node $v$ :
$d(v)=$ discovery time of node $v$. The time at which node $v$ is first reached. The time of preorder processing.
$f(v)=$ finish time of node $v$. The time at which node $v$ is exited for the last time. The time of postorder processing.

For a node $v$, active $(v)=$ time interval $d(v)$ to $f(v)$ (inclusive).
Example 1: In this particular depth-first search, the alphabeti-cally-first node is chosen, whenever an arbitrary choice is made.


Example 2: Depth-first search of the same digraph. The alpha-betically-last node is chosen, whenever an arbitrary choice is made.


Example 3: Depth-first search of the same digraph. In choosing among adjacent vertices not yet discovered, the alphabeticallyfirst vertex is chosen. However, we choose H as the starting vertex, and when the stack becomes empty (which didn't occur in Example 1 or 2 ), we choose E , then P , and then D as the next vertex to discover.


## Depth-first Search in Digraphs — edges classified

Edge $(u, v)$ is
a tree edge if $\operatorname{active}(u) \supset \operatorname{active}(v)$, and there is no vertex $x$ with $\operatorname{active}(u) \supset \operatorname{active}(x) \supset \operatorname{active}(v)$,
a forward edge if $\operatorname{active}(u) \supset \operatorname{active}(v)$, but $(u, v)$ is not a tree edge,
a back edge $\quad$ if $\operatorname{active}(u) \subset \operatorname{active}(v)$,
a cross edge if active $(u) \cap \operatorname{active}(v)=\varnothing$ (in which case $\operatorname{active}(v)$ entirely precedes $\operatorname{active}(u)$.
$\longrightarrow$ denotes a tree edge.
$--\rightarrow$ denotes a back edge.
-.... denotes a forward edge.

$\rightarrow \quad$ denotes a cross edge.

Here depth-first search of Example 1, with edges classified.


