

CS / MCS 401 Week #3–4 Exercises (Spring 2008)

Exercise 2.3–1 (page 36).

D. Prove by induction on n that the recurrence

$$C(n) = d + C(k) + C(n-k-1), \quad C(0) = 0,$$

where d is a constant and $0 \leq k \leq n-1$, has solution $C(n) = dn$.

E. Show how to merge three sorted arrays of length $n/3$ into a single sorted array of length n using approximately $5/3n$ comparisons, in the worst case.

F. Consider 3-way mergesort, in which an array of size n is divided into three subarrays of equal size, and the three subarrays are sorted by recursive calls to 3-way mergesort. The recurrence for the number of comparisons becomes $C(n) = 3C(n/3) + (5/3)n$, $C(1) = 0$, assuming n is a power of 3, say $n = 3^k$. Find an exact solution to this recurrence, when n is a power of 3. (Your final solution should involve only n , not k .) How does the number of comparisons performed by 3-way mergesort compare to the number performed by ordinary mergesort?

G. Find a solution to each recurrence below that is exact when n is a power of 2, and a good approximation otherwise.

- a) $C(n) = C(n/2) + 2n + 3$, $C(1) = 1$.
- b) $C(n) = 2C(n/2) + n \lg(n)$, $C(1) = 0$.
- c) $C(n) = 4C(n/2) + 3n^2$, $C(1) = 0$.