

Solutions to CS/MCS 401 Week #5 Exercises (Spring 2008)

Problem 4-4, parts (a), (c), (e), (h)

a) $T(n) = 3T(n/2) + n \lg(n)$.

In the Master Theorem, $a = 3$, $b = 2$, $E = \lg(3) \approx 1.59$, and $f(n) = n \lg(n)$. $f(n) = O(n^{E-\varepsilon})$, where we could take $\varepsilon = 0.1$. The Master Theorem (case 1) tells us $T(n) = \Theta(n^{\lg(3)}) \approx \Theta(n^{1.59})$.

c) $T(n) = 4T(n/2) + n^2 \text{sqrt}(n) = 4T(n/2) + n^{5/2}$.

In the Master Theorem, $a = 4$, $b = 2$, $E = \lg(4) = 2$, $n^E = n^2$, and $f(n) = n^{5/2}$. $f(n) = \Omega(n^{E+\varepsilon})$, where we could take $\varepsilon = 0.1$. $af(n/b) = 4(n/2)^{5/2} = 2^{-1/2}n$, so $af(n/b) \leq cf(n)$ where $c = 2^{-1/2} < 1$. The Master Theorem (case 3) tells us $T(n) = \Theta(f(n)) = \Theta(n^{5/2})$.

e) $T(n) = 2T(n/2) + n/\lg(n)$.

In the notation of the Master Theorem, $a = 2$, $b = 2$, $E = \lg(2) = 1$, $f(n) = n/\lg(n)$, and $n^E = n$. The Master Theorem does not apply, as $f(n)$ grows too rapidly for case (1) and not rapidly enough for case (2). But the extension to the Master Theorem in the handout tells us that $T(n) = \Theta(n \lg \lg(n))$.

h) $T(n) = T(n-1) + \lg(n)$

$$\begin{aligned} T(n) &= T(n-1) + \lg(n) \\ &= T(n-2) + \lg(n-1) + \lg(n) \\ &= T(n-3) + \lg(n-2) + \lg(n-1) + \lg(n) \\ &\quad \vdots \\ &= T(1) + \lg(2) + \lg(n-2) + \lg(n-1) + \lg(n) \\ &= T(1) + \lg(n!) \\ &= n \lg(n) - 1.44n + O(\lg(n)) = \Theta(n \lg(n)) \end{aligned}$$

Exercise H. Assume the recurrence

$$C(n) = C(0.8n) + C(0.5n) + C(0.2n) + n, \quad C(1) = 1$$

has a solution

$$C(n) = an^b + cn + d$$

for some real numbers a , b , c , and d . We ignore the fact that $0.2n$, $0.5n$, and $0.8n$ are not in general integers. Substituting the proposed solution into the recurrence, we obtain

$$\begin{aligned} an^b + cn + d &= a(0.2n)^b + c(0.2n) + d + \\ &\quad a(0.5n)^b + c(0.5n) + d + \\ &\quad a(0.8n)^b + c(0.8n) + d + n \end{aligned}$$

or

$$a(1 - 0.2^b - 0.5^b - 0.8^b)n^b + (c - 0.2c - 0.5c - 0.8c - 1)n - 2d = 0.$$

Since this must hold for all n , the coefficients of n^b , n , and 1 each must equal 0.

$$\begin{aligned} a(1 - 0.2^b - 0.5^b - 0.8^b) &= 0 \Rightarrow 1 - 0.2^b - 0.5^b - 0.8^b = 0 \quad (\text{since } a \neq 0), \\ c - 0.2c - 0.5c - 0.8c - 1 &= 0 \Rightarrow c = -2, \end{aligned}$$

$$2d = 0 \Rightarrow d = 0.$$

Now the condition $C(1) = 1$ implies $a + c + d = 1$, which together with the values of c and d above gives $a = 3$.

Let $f(x) = 1 - 0.2^x - 0.5^x - 0.8^x$. b is a root of $f(x)$. $f(x)$ is clearly an increasing function of x , since each of 0.2^x , 0.5^x , and 0.8^x decrease as x increases. $f(1) = -0.5$ and $f(2) = 0.07$. Thus $f(x)$ has a unique root (equal to b), which must lie between 1 and 2 (probably closer to 2), and it can be approximated by bisection or Newton's method. With an initial guess of 2, Newton's method with 4 iterations gives $b = 1.8267247$, or to two decimal places $b \approx 1.83$.

Note: The complete solution is $C(n) = 3n^{1.83} - 2n$. We saw earlier that division into 3 equal-sized subproblems of total size $1.5n$, with linear divide/combine time, led to a $\Theta(n^{1.59})$ time algorithm. Here the division into three unequal-sized subproblems with the same total size raises the time to $\Theta(n^{1.83})$.

Exercise I. The inversions in the array $\mathbf{a} = (41, 16, 74, 33, 66, 54)$ are:

$(41,16), (41,33), (74,33), (74,66), (74,54), (66,54)$.

The number of inversions is 6. Straight insertion sort would perform 6 comparisons in which it finds the elements out of order (1 for each inversion) and $n-1 = 5$ comparisons in which it finds the elements in order — a total of **11 comparisons**. Each comparison in which the elements are out of order is followed by an exchange, so there are **6 exchanges**.

Exercise J

- a) **Not a strict weak order.** Here $(x,y) \sim (u,v)$ when **not**($x < u$ and $y < v$) and **not**($u < x$ and $v < y$).
Note $(0,0) \sim (2,-1)$ and $(2,-1) \sim (1,2)$, but $(0,0) \not\sim (1,2)$, so \sim is not transitive.
- b) **A strict weak order.** The equivalence classes consist are the ellipses centered at the origin, with semi-major axis along the X-axis, and semi-minor axis equal to half the semi-major axis.
- c) **A strict weak order.** $(x,y) \sim (u,v)$ when $x - y = u - v$. For each real number α , there is an equivalence class consisting of those points (x,y) for which $x - y = \alpha$, or equivalently, $y = x - \alpha$. This class is the line with slope 1 and Y-intercept $-\alpha$.
- d) **Not a strict weak order.** Note $(x,y) \sim (u,v)$ when **not**($x < u - 1$) and **not**($u < x - 1$). Thus $(x,y) \sim (u,v)$ when $|x - u| \leq 1$. Note $(0,0) \sim (1,0)$ and $(1,0) \sim (2,0)$, but $(0,0) \not\sim (2,0)$, so \sim is not transitive.
- e) **A strict weak order.** For each integer i , there is an equivalence class consisting of $\{r \mid r \text{ is a real number with } i \leq r < i+1\}$.
- f) **A strict weak order.** For each real number α in $[0,1)$, there is an equivalence class consisting of $\{\alpha + i \mid i \text{ is an integer}\}$.
- g) **Not a strict weak order.** Note $a \sim b$ if neither a nor b is a proper divisor of the other. So $2 \sim 3$ and $3 \sim 4$, but $2 \not\sim 4$. This means \sim is not transitive.