

The Algorithm for Finding the k^{th} Smallest in Expected Linear Time (Non-Recursive)

Let T be the type of the elements in array a . As the algorithm runs, at all times the k^{th} smallest element of a lies between positions $left$ and $right$, inclusive. The algorithm terminates when $left == right$, at which time both must equal k .

```
T select( T[] a, Integer k)
    left = 1;
    right = a.length;
    while ( left < right )
        q = partition( a, left, right);
        if ( k < q )
            right = q - 1;
        else if ( k > q )
            left = q + 1;
        else
            left = right = q;
    return a[q];
```

A Recursive Implementation of The Algorithm for Finding the k^{th} Smallest in Expected Linear Time

The recursive version of *select()* below is initiated by invoking *select(a, k, 1, a.length)*. It assumes the array *a* is passed by reference and the integer arguments by value, as in C, C++, or Java.

When *select(a, k, left, right)* is called, the k^{th} smallest element of the entire array *a* must lie between positions *left* and *right*, inclusive. *select()* completes the computation of the k^{th} smallest element and returns it.

```
T select( T[] a, Integer k, Integer left, Integer right)
    if ( left < right )
        q = partition( a, left, right);
        if ( k < q )
            return select( a, k, left, q - 1);
        else if ( k > q )
            return select( a, k, q + 1, right);
        else
            return a[q];
    else
        return a[left];
```